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# **Graph Theory & Social Network Analysis**

#### **Outline**

#### **Network properties**

- *Adjacency matrices*
- *Paths, shortest paths*
- *Network diameter*

#### **Node properties**

- *Degree*
- *Centrality*
- *Clustering coefficient*



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# Network properties

Adjacency matrices Paths, shortest paths Network diameter 

#### **Graph G = (V,E)**

- $-V =$  set of nodes
- $E = set of edges$





#### **Networks as graphs**





An **undirected graph** is one in which edges have **no orientation.** The edge (a, b) is identical to the edge (b, a).

**a directed graph** is a graph, or set of nodes connected by edges, where the edges *have a* direction associated with them.

# **Adjacency matrix**

- **Representing edges (who is** adjacent to whom) in a matrix
	- $-$  A<sub>ij</sub> = 1 if node *i* has an edge to node *j* **= 0** if node *i* does not have an edge to j
	- $A_{ii}$  = 0 unless the network has self-loops
	- $-$  **A**<sub>ii</sub> = **A**<sub>ii</sub> if the network is undirected, or if *i* and *j* share a reciprocated edge



#### **Adjacency matrix example**



# **Walks, Paths, Cycles, and Geodesics**

• Walk from i to ik: a sequence of nodes (i<sub>1,</sub>i<sub>2,...</sub> ik) and a sequence of links (i<sub>1</sub>i<sub>2</sub>, i<sub>2</sub>i<sub>3</sub>, ..., i<sub>k-1</sub>i<sub>k</sub>) such that i<sub>k-1</sub>i<sub>k</sub> in E for each k

• **Path:** a walk  $(i_1, i_2, \ldots, i_k)$  with each node  $i_k$  is distinct

• **Cycle:** a walk where  $i_1 = i_k$ 

• **Geodesic:** a shortest path between two nodes

#### **Network Diameter**

- Diameter = the *longest shortest path* in the network
	- $-$  Represents a worst-case scenario in network size
	- Left example (undirected network): diameter=?
	- Right example (directed network): diameter=?



#### **Diameter scenarios**





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# Node properties

Degree **Centrality** Clustering Coefficient

## **Degree of Nodes**

**Degree:** number of edges incident on a node **abula container to the degree=5** 

**Two different degree types in directed networks** 

**1. Indegree** : how many directed edges (arcs) terminate at a node

**2. Outdegree**: how many directed edges (arcs) originate at a node





## **Node degree from matrix values**



**example:** outdegree for node 3 is 2, which we obtain by summing the number of non-zero entries in the 3<sup>rd</sup> *row* 

**example:** the indegree for node 3 is 1, which we obtain by summing the number of nonzero entries in the 3rd column

# **Degree distribution**

- **Degree distribution**: A frequency count of the occurrence of each degree in a network
- Degree distributions are far from normal in most real-world networks (hubs)

**Example:** In the figure we witness many nodes with very small degree and few nodes with high degree, implying the presence of a **hub** or fat tail.



**Centrality:** Captures the idea of how central a node is in the network

#### *Can be categorized into four main types*

- **1. Degree Centrality:** Shows how connected a node is
- **2. Betweenness Centrality:** Shows how important a node is in terms of connecting other nodes
- **3. Closeness Centrality:** Shows how easily a node can reach other nodes (i.e. how close the node is to the center of the network)
- **4. Eigenvector / Bonacich Centrality:** Show how much a node is connected to other important nodes in the network.

#### In each of the following networks, X has higher centrality than Y according to a particular measure



# **Degree Centrality**

Nodes with more friends are more important



**Assumption:** the connections that your friends have don't matter, it is what they can do directly that does

# **Normalization of Degree Centrality**

- Divide degree by the maximum possible (N-1)
- Normalized Degree Centrality ranges from 0 to 1
- Allows comparisons between networks of different sizes



## **Centralization: skew in distribution**

- **Centrality refers to an individual node but there is a need to** capture the inequality in the distribution of centralities characterizing the entire network.
- Using Freeman's general formula for **centralization** we can capture the inequality of degree between the nodes of the network:



#### **Degree centralization examples**



$$
C_D=0.167
$$

 $\overline{2}$ 

#### In what ways does degree fail to capture centrality in the following graphs?



**Brokerage not captured by degree !** 

How many pairs of individuals would have to go through you in order to reach one another in the minimum number of hops?

# **X Y**

#### **Betweenness Centrality**

$$
C_B(i) = \sum_{j < k} g_{jk}(i) / g_{jk}
$$
\n
$$
\cdot g_{jk} = \text{the number of shortest paths connecting } jk
$$
\n
$$
\cdot g_{jk}(i) = \text{the number that actor } i \text{ is on.}
$$

#### **Usually normalized by:**

$$
C'_{B}(i) = C_{B}(i) \sqrt{[(n-1)(n-2)/2]}
$$
  
number of pairs of vertices  
excluding the vertex itself

#### **Betweenness on toy networks**

# **non-normalized version**



- A lies between no two other vertices
- B lies between A and 3 other vertices: C, D, and E
- C lies between 4 pairs of vertices  $(A,D)$ ,  $(A,E)$ ,  $(B,D)$ ,  $(B,E)$
- note that there are no alternate paths for these pairs to take, so C gets full credit

## **non-normalized version**



#### **Betweenness on toy networks**

# **non-normalized version**



- why do C and D each have betweenness 1?
- They are both on shortest paths for pairs  $(A, E)$ , and  $(B, E)$ , and so must share credit:
	- $\frac{1}{2} + \frac{1}{2} = 1$

#### **Betweenness centrality in directed networks**

• We now consider the fraction of all directed paths between any two vertices that pass through a node



Only modification: when normalizing, we have  $(N-1)*(N-2)$  instead of  $(N-1)*(N-2)/2$ , because we have twice as many ordered pairs as unordered pairs

$$
C'_{B}(i) = C_{B}(i)/[(N-1)(N-2)]
$$

# **Closeness Centrality**

- What if it's not so important to have many direct friends?
- Or be "between" others
- But one still wants to be in the "middle" of things, not too far from the center



**Closeness** is based on the length of the average shortest path between a node and all other nodes in the network

**Closeness Centrality:** 

$$
C_c(i) = \left[\sum_{j=1}^N d(i,j)\right]^{-1}
$$

**Normalized Closeness Centrality** 

 $C_C^{\dagger}(i) = (C_C(i))/(N-1)$ 

#### **Closeness Centrality Toy Example**



$$
C_c^{\dagger}(A) = \left[ \frac{\sum_{j=1}^{N} d(A,j)}{N-1} \right]^{-1} = \left[ \frac{1+2+3+4}{4} \right]^{-1} = \left[ \frac{10}{4} \right]^{-1} = 0.4
$$

#### **Closeness Centrality Other Examples**





# **Directed closeness centrality**

- **Choose a direction** 
	- $-$  **in-closeness** (e.g. prestige in citation networks)
	- **out-closeness**
- Usually consider only vertices from which the node in question can be reached



#### How central you are depends on how central your **neighbors** are



While the **degree** for node A in a social network measures how  $m$  any ties  $A$  has, the **eigenvector centrality** of node A is measured based on **how** *many ties A's connections have.* 

#### **Eigenvector centrality**

#### How central you are depends on how central your **neighbors** are



The centrality score  $c_{(i)}$  of each node i is proportional to its neighbors' scores

$$
C(i) = A_{ji}C(j) + A_{ki}C(k) + A_{li}C(l)
$$

#### **Bonacich eigenvector centrality**

• The **Bonacich Centrality** measure is also based on the premise that a node's importance is determined by *how important its neighbors* are.

$$
c_i(\beta) = \sum_j (\alpha + \beta c_j) A_{ji}
$$

- $\alpha$  is a normalization constant
- $\beta$  determines how important the centrality of your neighbors is
- $\cdot$  **A** is the adjacency matrix (can be weighted)
- This notion is central to citation rankings and things like Google page rankings.

# **Bonacich Power Centrality: attenuation factor b**

#### **small**  $\beta \rightarrow$  high attenuation

- only your immediate friends matter, and their importance is factored in *only a bit*
- high  $\beta \rightarrow$  low attenuation
	- global network structure matters (your friends, your friends' of friends *etc.)*
- $\beta$  = 0 yields simple degree centrality
- If β > 0, nodes have higher centrality when they have edges to other central nodes.
- If β < 0, nodes have higher centrality when they have edges to less central nodes.

#### **Bonacich Power Centrality: examples**



#### **Why does the middle node have lower centrality than its neighbors when β is negative?**

#### **Example Centrality Measures**





**(Local) clustering coefficient for a node is the probability that two** randomly selected friends of a node are friends with each other



Fraction of the friends of a node that are friends with each other (i.e., connected)

# **Clustering Coefficient**





#### \* Ranges from 0 to 1